

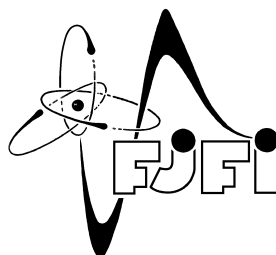
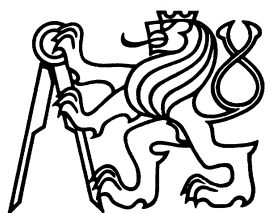
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**Morphisms preserving the set of words  
coding three interval exchange**

**Morfismy zachovávající slova kódující  
výměnu tří intervalů**

**VÝZKUMNÝ ÚKOL**

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Tomáš Hejda

## **Prohlášení**

Prohlašuji, že jsem předloženou práci vypracoval samostatně a že jsem uvedl veškerou použitou literaturu.

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**Title:** Morphisms preserving the set of words coding three interval exchange

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### **Abstract**

Any amicable pair  $\varphi, \psi$  of Sturmian morphisms enables a construction of a ternary morphism  $\eta$  which preserves the set of infinite words coding 3-interval exchange. We determine the number of amicable pairs with the same incidence matrix in  $SL(2, \mathbb{N})$  and we study incidence matrices associated with the corresponding ternary morphisms  $\eta$ .

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## 1 Introduction

*Sturmian words* are well-described objects in combinatorics on words. They can be defined in several equivalent ways [5], e.g. as words coding a two-interval exchange transformation with irrational ratio of lengths of the intervals. Morphisms preserving the set of Sturmian words are called *Sturmian* and they form a monoid generated by three of its elements (see [6, 12]). Let us denote this monoid by  $\mathcal{M}_{\text{Sturm}}$ .

In this paper, we consider morphisms preserving the set of words coding a three-interval exchange transformation with permutation  $(3, 2, 1)$ , the so-called *3iet words*. We call these morphisms *3iet-preserving*. Monoid of these morphisms, denoted by  $\mathcal{M}_{\text{3iet}}$ , is not fully described. It is shown (see [9]) that the monoid  $\mathcal{M}_{\text{3iet}}$  is not finitely generated. Recently, in [2], pairs of amicable Sturmian morphisms were defined. The authors used this notion to describe morphisms that have as a fixed point a non-degenerate 3iet word, i.e. word with complexity  $\mathcal{C}(n) = 2n + 1$ . Using the operation of “ternarization”, we can assign a morphism  $\eta = \text{ter}(\varphi, \psi)$  over a ternary alphabet to a pair of amicable Sturmian morphisms. We show that such  $\eta$  is a 3iet-preserving morphism. Moreover, we show that the set

$$\mathcal{M}_{\text{ter}} = \{ \text{ter}(\varphi, \psi) \mid \varphi, \psi \text{ amicable morphisms} \} \quad (1)$$

is a monoid, but it does not cover the whole monoid  $\mathcal{M}_{\text{3iet}}$ .

We also study the incidence matrices of morphisms  $\eta \in \mathcal{M}_{\text{ter}}$ . From the definition of amicable Sturmian morphisms  $\varphi, \psi$  we can derive that  $\varphi$  and  $\psi$  have the same incidence matrix  $\mathbf{A} \in \mathbb{N}^{2 \times 2}$ , where  $\det \mathbf{A} = \pm 1$ . As shown in [13], for every matrix  $\mathbf{A} = \begin{pmatrix} p_0 & q_0 \\ p_1 & q_1 \end{pmatrix}$  with  $\det \mathbf{A} = \pm 1$ , there exist  $p_0 + p_1 + q_0 + q_1 - 1$  Sturmian morphisms. We will show the following theorem concerning the number of pairs of amicable Sturmian morphisms with a given matrix.

**Theorem 1.** *Let  $\mathbf{A} = \begin{pmatrix} p_0 & q_0 \\ p_1 & q_1 \end{pmatrix} \in \mathbb{N}^{2 \times 2}$  be a matrix with  $\det \mathbf{A} = \pm 1$ . Then there exist exactly*

$$m(\|\mathbf{A}\| - 1) + \frac{m}{2}(\det \mathbf{A} - m) \quad (2)$$

*pairs of amicable Sturmian morphisms with incidence matrix  $\mathbf{A}$ , where  $m = \min\{p_0 + p_1, q_0 + q_1\}$  and  $\|\mathbf{A}\| = p_0 + p_1 + q_0 + q_1$ .*

Moreover, for a given matrix  $\mathbf{A}$ , we will describe all matrices  $\mathbf{B} \in \mathbb{N}^{3 \times 3}$  such that  $\mathbf{B}$  is an incidence matrix of  $\eta = \text{ter}(\varphi, \psi)$  for amicable Sturmian morphisms  $\varphi, \psi$  with incidence matrix  $\mathbf{A}$ .

## 2 Preliminaries

### 2.1 Words over a finite alphabet

Besides the infinite words, we consider *finite words* over the alphabet  $\mathcal{A}$ . We write  $w = w_0w_1 \cdots w_{n-1}$ , where  $w_i \in \mathcal{A}$  for all  $i \in \mathbb{N}$ ,  $i < n$ . We denote by  $|w|$  the length  $n$  of the finite word  $w$ . We denote by  $|w|_a$  the number of occurrences of a letter  $a \in \mathcal{A}$  in the word  $w$ . The set of all finite words on the alphabet  $\mathcal{A}$  including the empty word  $\epsilon$  is denoted by  $\mathcal{A}^*$ . The set  $\mathcal{A}^*$  with the operation of concatenation is a monoid. On the set  $\mathcal{A}^*$  we define a relation of *conjugation*:  $w \sim w'$ , if there exists  $v \in \mathcal{A}^*$  such that  $wv = vw'$ . A *morphism* from  $\mathcal{A}^*$  to  $\mathcal{B}^*$  is a mapping  $\varphi : \mathcal{A}^* \rightarrow \mathcal{B}^*$  such that  $\varphi(vw) = \varphi(v)\varphi(w)$  for all  $v, w \in \mathcal{A}^*$ . It is clear that a morphism is well defined by images of letters  $\varphi(a)$  for all  $a \in \mathcal{A}$ . If  $\mathcal{A} = \mathcal{B}$ , then  $\varphi$  is called a *morphism over  $\mathcal{A}$* .

The set of *infinite words* over the alphabet  $\mathcal{A}$  is denoted by  $\mathcal{A}^{\mathbb{N}}$ . The action of a morphism can be naturally extended to an infinite word  $(u_i)_{i \in \mathbb{N}}$  putting  $\varphi(u) = \varphi(u_0)\varphi(u_1)\varphi(u_2) \cdots$ . If an infinite word  $u \in \mathcal{A}^{\mathbb{N}}$  satisfies  $\varphi(u) = u$ , we call it a *fixed point* of the morphism  $\varphi$  over  $\mathcal{A}$ .

To a morphism  $\varphi$  over  $\mathcal{A}$  we assign an *incidence matrix*  $\mathbf{M}_\varphi$  defined by  $(\mathbf{M}_\varphi)_{ab} = |\varphi(a)|_b$  for all  $a, b \in \mathcal{A}$ . To a finite word  $v \in \mathcal{A}^*$  we assign a *Parikh vector*  $\mathbf{M}_v$  defined by  $(\mathbf{M}_v)_b = |v|_b$  for all  $b \in \mathcal{A}$ .

The *language* of an infinite word  $u$  is the set of all its factors. Let us recall that a finite word  $w \in \mathcal{A}^*$  is a *factor* of  $u = (u_i)_{i \in \mathbb{N}}$ , if there exist indices  $n, j \in \mathbb{N}$  such that  $w = u_nu_{n+1} \cdots u_{n+j-1}$ . The language of an infinite word is denoted by  $\mathcal{L}(u)$ .

It is known that the language of neither Sturmian nor 3iet word depends on the point  $x_0 \in [0, 1)$ , the orbit of which the infinite word codes. It depends only on slope  $\varepsilon$  or parameters  $\alpha, \beta$ .

The *(factor) complexity* of an infinite word  $u$  is a mapping  $\mathcal{C}_u : \mathbb{N} \rightarrow \mathbb{N}$ , which returns the number of factors of  $u$  of the length  $n$ , thus  $\mathcal{C}_u(n) = \#\{w \in \mathcal{L}(u) \mid |w| = n\}$ . It is easy to see that a word  $u$  is periodic if and only if there exists  $n_0 \in \mathbb{N}$  such that  $\mathcal{C}_u(n_0) \leq n_0$ .

### 2.2 Interval exchange

We consider Sturmian words, i.e. aperiodic words given by exchange of 2 intervals with permutation  $(2, 1)$ , and words given by exchange of 3 intervals with permutation  $(3, 2, 1)$ . Let us recall that general  $r$ -interval exchange transformations were introduced already in [11].

**Two-interval exchange.** The 2-interval exchange transformation  $S$  is a mapping  $S : [0, 1) \rightarrow [0, 1)$ . It is determined by its slope  $\varepsilon \in [0, 1]$  and is

given by

$$Sx = \begin{cases} x + 1 - \varepsilon & \text{if } x \in [0, \varepsilon) \\ x - \varepsilon & \text{if } x \in [\varepsilon, 1). \end{cases}$$

The orbit of a point  $x_0 \in [0, 1)$  with respect to the transformation  $S$ , i.e. the sequence  $x_0, Sx_0, S^2x_0, \dots$  can be coded by an infinite word  $u = (u_i)_{i=0}^{\infty}$  on the binary alphabet  $\{0, 1\}$ . The infinite word is given by

$$u_i = \begin{cases} 0 & \text{if } S^i x_0 \in [0, \varepsilon), \\ 1 & \text{if } S^i x_0 \in [\varepsilon, 1). \end{cases} \quad (3)$$

It is a well-known fact that for an irrational  $\varepsilon$ , the word  $u$  is Sturmian. Using the same construction on the partition of the interval  $(0, 1]$  into  $(0, \varepsilon] \cup (\varepsilon, 1]$ , we again obtain a Sturmian word. On the other hand, every Sturmian word can be obtained by one of the above two constructions. The set of Sturmian words will be denoted by  $\mathcal{W}_{\text{Sturm}}$ .

In [12], the authors show that Sturmian words are the aperiodic words with minimal complexity, i.e.  $\mathcal{C}_u(n) = n + 1$  for all  $u \in \mathcal{W}_{\text{Sturm}}$  and  $n \in \mathbb{N}$ . We can see that

$$S^i x_0 = \{x_0 - i\varepsilon\} \quad \text{for all } x_0 \in [0, 1), \quad (4)$$

where  $\{x\} = x - \lfloor x \rfloor$  denotes the *fractional part* of a number  $x \in \mathbb{R}$ . Then  $u_i = \lfloor x_0 - i\varepsilon \rfloor - \lfloor x_0 - (i+1)\varepsilon \rfloor$ , which is exactly the formula how [12] define mechanical words.

We will use another fact about two-interval exchange. Let  $\varphi \in \mathcal{M}_{\text{Sturm}}$  be a Sturmian morphism. Then the word  $\varphi(a)$  for  $a \in \{0, 1\}$  codes two-interval exchange with the slope  $\frac{|v|_0}{|v|}$ . We should see this from the Lemma 2.1.15 in [12]. The word  $a^k$  is a factor of some Sturmian word, hence the word  $\varphi(a)^k$  is balanced for any  $k \in \mathbb{N}$ , which means that the infinite word  $u = \varphi(a)^\omega = \varphi(a)\varphi(a)\varphi(a)\dots$  is balanced and periodic, thus it is rational mechanical. In our terms, this means that it codes a rational 2-interval exchange; it is as well shown there that the slope of the transformation is exactly  $\frac{|v|_0}{|v|}$ .

**Three-interval exchange.** The 3-interval exchange transformation  $T$  is determined by two parameters  $\alpha, \beta \in (0, 1)$  satisfying  $\alpha + \beta < 1$ . Using parameters  $\alpha, \beta$  and  $\gamma = 1 - \alpha - \beta$  we partition the interval  $[0, 1)$  into  $I_A = [0, \alpha)$ ,  $I_B = [\alpha, \alpha + \beta)$  and  $I_C = [\alpha + \beta, 1)$ . The mapping  $T$  is given by

$$Tx = \begin{cases} x + \beta + \gamma & \text{if } x \in I_A, \\ x - \alpha + \gamma & \text{if } x \in I_B, \\ x - \alpha - \beta & \text{if } x \in I_C. \end{cases}$$

The orbit of a point  $x_0 \in [0, 1)$  with respect to the transformation  $T$  is coded by a word  $u = (u_i)_{i=0}^{\infty}$  over the ternary alphabet  $\{A, B, C\}$ :

$$u_i = X \quad \text{if} \quad T^i x_0 \in I_X.$$

Similarly to the case of 2-interval exchange transformation, we can define the exchange of 3 intervals using the partition  $(0, 1] = (0, \alpha] \cup (\alpha, \alpha + \beta] \cup (\alpha + \beta, 1]$ . If  $\frac{1-\alpha}{1+\beta}$  is irrational, the infinite word  $u$  is aperiodic, and we call it a *3iet word*; the set of these words is denoted by  $\mathcal{W}_{3iet}$ . For combinatorial properties of 3iet words, see [8].

Aperiodic words coding 3-interval exchange transformations, called here 3iet words, have the complexity  $\mathcal{C}_u(n) \leq 2n + 1$  for all  $n \in \mathbb{N}$ . If a 3iet word  $u \in \mathcal{W}_{3iet}$  satisfies  $\mathcal{C}_u(n) = 2n + 1$  for all  $n \in \mathbb{N}$ , we call it a *non-degenerate* 3iet word; otherwise we call it a *degenerate* 3iet word and it is a quasi-Sturmian word (see [7]).

### 2.3 Standard pairs and standard morphisms

In [13], the notion of standard pairs is introduced. If we define two operators on pairs of words  $L, R : \mathcal{A}^* \times \mathcal{A}^* \rightarrow \mathcal{A}^* \times \mathcal{A}^*$  as

$$L(x, y) = (x, xy), \quad R(x, y) = (yx, y),$$

we say that a pair  $(x, y)$  is a *standard pair*, if it can be obtained from the pair  $(0, 1)$  by applying the operators  $L$  and  $R$  finitely many times.

We say that a binary morphism  $\varphi$  is *standard*, if there exists a standard pair  $(x, y)$  such that

$$\begin{array}{ccc} \varphi(0) = x, & & \varphi(0) = y, \\ \varphi(1) = y, & \text{or} & \varphi(1) = x. \end{array}$$

The authors of [13] show the close connection between the standard morphisms and all the Sturmian morphisms:

1. Every standard morphism is Sturmian.
2. For every matrix  $\mathbf{A} \in \mathbb{N}^{2 \times 2}$  with  $\det \mathbf{A} = \pm 1$ , there exists exactly one standard morphism  $\varphi$  with incidence matrix  $\mathbf{M}_\varphi = \mathbf{A}$ .
3. Every Sturmian morphism  $\psi \in \mathcal{M}_{\text{Sturm}}$  is a right conjugate to some standard morphism  $\varphi$ . Let us recall that a morphism  $\psi$  over  $\mathcal{A}$  is a *right conjugate* to  $\varphi$ , if there exists a finite word  $v \in \mathcal{A}^*$  such that

$$\psi(a)v = v\varphi(a) \quad \text{for all letters } a \in \mathcal{A}.$$



## 2.4 Amicable words and morphisms

In the article [4], authors show the close connection between 3iet and Sturmian words using two morphisms  $\sigma_{01}, \sigma_{10} : \{A, B, C\}^* \rightarrow \{0, 1\}^*$  given by

$$\begin{aligned} \sigma_{01}(A) &= 0, & \sigma_{10}(A) &= 0, \\ \sigma_{01}(B) &= 01, & \sigma_{10}(B) &= 10, \\ \sigma_{01}(C) &= 1, & \sigma_{10}(C) &= 1. \end{aligned}$$

In [4], the following theorem is proved.

**Theorem 2.** *An infinite ternary word  $u \in \{A, B, C\}^{\mathbb{N}}$  is a 3iet word if and only if the words  $\sigma_{01}(u)$  and  $\sigma_{10}(u)$  are Sturmian.*

This theorem motivated the authors of [1] to introduce the relation of amicability of words.

**Definition 3.** Let  $w, w' \in \{0, 1\}^*$ , let  $b \in \mathbb{N}$ . We say that  $w$  is *b-amicable* to  $w'$ , if there exists a factor  $v \in \{A, B, C\}^*$  of some 3iet word such that

$$w = \sigma_{01}(v), \quad w' = \sigma_{10}(v) \quad \text{and} \quad |v|_B = b.$$

We say that  $w$  is *amicable* to  $w'$ , if  $w$  is *b-amicable* to  $w'$  for some  $b \in \mathbb{N}$ , and we denote it by  $w \propto w'$ .

The ternary word  $v$  is called a *ternarization* of  $w$  and  $w'$ , and we write  $v = \text{ter}(w, w')$ .

It is easy to see that if  $w \propto w'$ , then they are factors of the same Sturmian word and their Parikh vectors coincide.

In [1], the notion of “amicable words” plays a crucial role in enumeration of words with length  $n$  occurring in a 3iet word. In [2], the authors investigate ternary morphisms that have a non-degenerate 3iet fixed point using the following notion of amicability of two Sturmian morphisms.

**Definition 4.** Let  $\varphi, \psi$  be Sturmian morphisms over the alphabet  $\{0, 1\}$ . We say that  $\varphi$  is *amicable* to  $\psi$ , if

$$\begin{aligned} \varphi(0) &\propto \psi(0), \\ \varphi(01) &\propto \psi(10) \\ \text{and } \varphi(1) &\propto \psi(1). \end{aligned}$$

We denote this relation by  $\varphi \propto \psi$ . The morphism  $\eta$  over the ternary alphabet  $\{A, B, C\}$ , given by

$$\begin{aligned} \eta(A) &= \text{ter}(\varphi(0), \psi(0)), \\ \eta(B) &= \text{ter}(\varphi(01), \psi(10)), \\ \eta(C) &= \text{ter}(\varphi(1), \psi(1)) \end{aligned}$$

is called the *ternarization* of morphisms  $\varphi$  and  $\psi$ , and is denoted by  $\eta = \text{ter}(\varphi, \psi)$ . Set of these  $\eta$  is denoted by  $\mathcal{M}_{\text{ter}}$ .

The article [2] states the following theorem:

**Theorem 5.** *Let  $\eta$  be a ternary morphism with non-degenerate 3iet fixed point. Then  $\eta \in \mathcal{M}_{\text{ter}}$  or  $\eta^2 \in \mathcal{M}_{\text{ter}}$ .*

### 3 Main results

#### 3.1 Globally 3iet-preserving morphisms

Analogously to the terminology introduced for Sturmian words and morphisms in [6], the ternarization  $\eta$ , having a 3iet fixed point, is *locally 3iet-preserving*, i.e. there exists  $u \in \mathcal{W}_{3\text{iet}}$  such that  $\eta(u) \in \mathcal{W}_{3\text{iet}}$ . We now prove a partial result about (*globally*) 3iet-preserving morphisms, i.e. ternary morphisms  $\eta$  such that

$$\eta(u) \in \mathcal{W}_{3\text{iet}} \quad \text{for all } u \in \mathcal{W}_{3\text{iet}}.$$

**Proposition 6.** *Let  $\eta = \text{ter}(\varphi, \psi)$  for amicable Sturmian morphisms  $\varphi \propto \psi$ . Then  $\eta$  is a globally 3iet-preserving morphism.*

*Proof.* Directly from definitions we see that

$$\begin{aligned} \sigma_{01}\eta(A) &= \varphi(0), & \sigma_{01}\eta(B) &= \varphi(01), & \sigma_{01}\eta(C) &= \varphi(1), \\ \sigma_{10}\eta(A) &= \psi(0), & \sigma_{10}\eta(B) &= \psi(10), & \sigma_{10}\eta(C) &= \psi(1). \end{aligned}$$

Therefore

$$\sigma_{01}\eta(v) = \varphi\sigma_{01}(v) \quad \text{and} \quad \sigma_{10}\eta(v) = \psi\sigma_{10}(v) \quad (5)$$

for any factor  $v$  of a 3iet word  $u \in \mathcal{W}_{3\text{iet}}$ . According to Theorem 2 we get that  $\sigma_{01}(u)$  and  $\sigma_{10}(u)$  are Sturmian words, and since  $\varphi$  and  $\psi$  are Sturmian morphisms, we obtain that  $\sigma_{01}\eta(u)$  and  $\sigma_{10}\eta(u)$  are Sturmian words as well, which means, according to the same theorem, that the word  $\eta(u)$  is 3iet.  $\square$

**Proposition 7.** *Let  $\varphi_i \propto \psi_i$  be Sturmian morphisms, for  $i = 1, 2$ . Then*

$$\text{ter}(\varphi_1, \psi_1) \circ \text{ter}(\varphi_2, \psi_2) = \text{ter}(\varphi_1 \circ \varphi_2, \psi_1 \circ \psi_2).$$

*Proof.* It can be shown that the relation of amicability is preserved by composition of morphisms. More precisely  $\varphi_1\varphi_2 \propto \psi_1\psi_2$ . Denote  $\eta_1 = \text{ter}(\varphi_1, \psi_1)$ ,  $\eta_2 = \text{ter}(\varphi_2, \psi_2)$ . Using the relation (5), we see that for all  $v \in \{A, B, C\}^*$

$$\begin{aligned} \sigma_{01}\eta_1\eta_2(v) &= \varphi_1\sigma_{01}\eta_2(v) = \varphi_1\varphi_2\sigma_{01}(v) \\ \text{and } \sigma_{10}\eta_1\eta_2(v) &= \psi_1\sigma_{10}\eta_2(v) = \psi_1\psi_2\sigma_{10}(v). \end{aligned}$$

But this means that  $\eta_1\eta_2 = \text{ter}(\varphi_1\varphi_2, \psi_1\psi_2)$ .  $\square$

As a consequence of previous two propositions, we can state the following theorem.

**Theorem 8.** *The set  $\mathcal{M}_{\text{ter}}$  of all ternarizations of amicable Sturmian morphisms with the operation of composition of morphisms is a sub-monoid of the monoid  $\mathcal{M}_{\text{3iet}}$  of all globally 3iet-preserving morphisms.*

Unfortunately,  $\mathcal{M}_{\text{ter}} \subsetneq \mathcal{M}_{\text{3iet}}$ . Consider for example the morphism

$$\eta(A) = B, \quad \eta(B) = CAC, \quad \eta(C) = C. \quad (6)$$

As shown in [9], this morphism is 3iet-preserving, but it can be easily verified that it is not a ternarization of any pair of Sturmian morphisms, using the following statement.

**Proposition 9.** *A ternary morphism  $\eta$  is a ternarization, i.e.  $\eta \in \mathcal{M}_{\text{ter}}$ , if and only if it satisfies*

$$\sigma_{01}\eta(B) = \sigma_{01}\eta(AC) \quad \text{and} \quad \sigma_{10}\eta(B) = \sigma_{10}\eta(CA).$$

*Proof.* The implication ( $\Rightarrow$ ). Suppose  $\eta = \text{ter}(\varphi, \psi)$ . According to (5) we get

$$\begin{aligned} \sigma_{01}\eta(B) &= \varphi\sigma_{01}(B) = \varphi(01) = \varphi\sigma_{01}(AC) = \sigma_{01}\eta(AC), \\ \sigma_{10}\eta(B) &= \psi\sigma_{10}(B) = \psi(10) = \psi\sigma_{10}(CA) = \sigma_{10}\eta(CA). \end{aligned}$$

The implication ( $\Leftarrow$ ). Define morphisms  $\varphi, \psi$  as

$$\begin{aligned} \varphi(0) &= \sigma_{01}\eta(A), & \psi(0) &= \sigma_{10}\eta(A), \\ \varphi(1) &= \sigma_{01}\eta(C), & \psi(1) &= \sigma_{10}\eta(C). \end{aligned}$$

Immediately we get  $\text{ter}(\varphi(0), \psi(0)) = \eta(A)$  and  $\text{ter}(\varphi(1), \psi(1)) = \eta(C)$ . The words  $\varphi(01)$  and  $\psi(10)$  satisfy  $\varphi(01) = \sigma_{01}\eta(AC) = \sigma_{01}\eta(B)$  and  $\psi(10) = \sigma_{10}\eta(CA) = \sigma_{10}\eta(B)$ , which means that  $\text{ter}(\varphi(01), \psi(10)) = \eta(B)$ .  $\square$

For the morphism (6), we get  $\sigma_{01}\eta(B) = 010 \neq 001 = \sigma_{01}\eta(AC)$ . Another even simpler example of a 3iet-preserving morphism that is not a ternarization is the morphism interchanging the letters  $A$  and  $C$ .

### 3.2 Pairs of amicable Sturmian morphisms

Now, our goal will be to determine the number of amicable pairs of morphisms with incidence matrix  $\mathbf{A}$  of  $\det \mathbf{A} = \pm 1$ . We will use the notion of  $b$ -amicable morphisms.

**Definition 10.** Let  $\varphi$  and  $\psi$  be binary morphisms and let  $b \in \mathbb{N}$ . We say that  $\varphi$  is  $b$ -amicable to  $\psi$ , if  $\varphi$  is amicable to  $\psi$  and the number of occurrences of  $B$  in  $\text{ter}(\varphi(01), \psi(10))$  is  $b$ .

We now determine the numbers of pairs of  $b$ -amicable Sturmian morphisms. The following proposition and the Theorem 1 were already proven in [10]. We provide a more straightforward proof.

**Proposition 11.** *Let  $\mathbf{A} = \begin{pmatrix} p_0 & q_0 \\ p_1 & q_1 \end{pmatrix} \in \mathbb{N}^{2 \times 2}$  be a matrix with  $\det \mathbf{A} = \pm 1$  and  $b \in \mathbb{N}$ . Put  $p = p_0 + p_1$ ,  $q = q_0 + q_1$ . Then the number  $c_{\mathbf{A}}(b)$  of pairs of  $b$ -amicable morphisms with matrix  $\mathbf{A}$  is equal to*

$$c_{\mathbf{A}}(b) = \begin{cases} \|\mathbf{A}\| - b & \text{if } \det \mathbf{A} = +1 \text{ and } 1 \leq b \leq \min\{p, q\}, \\ \|\mathbf{A}\| - b - 2 & \text{if } \det \mathbf{A} = -1 \text{ and } 0 \leq b \leq \min\{p, q\} - 1, \\ 0 & \text{otherwise.} \end{cases}$$

First, let us state the following lemma.

**Lemma 1.** *Let  $\mathbf{A} = \begin{pmatrix} p_0 & q_0 \\ p_1 & q_1 \end{pmatrix} \in \mathbb{N}^{2 \times 2}$  be a matrix with  $\det \mathbf{A} = \pm 1$  and  $b \in \mathbb{N}$ . Put  $p = p_0 + p_1$ ,  $q = q_0 + q_1$  and  $N = p + q$ . Let  $S$  be a two-interval exchange with the slope  $p/(p + q)$ . Let  $w^{(k)}$  be a word of the length  $N$  that codes  $S$  with the start point  $k/N$ , for  $k \in \{0, \dots, N - 1\}$ .*

*Then  $w^{(k)}$  is  $b$ -amicable to  $w^{(\bar{k})}$  if and only if  $0 \leq b \leq \min\{p, q\}$  and  $\bar{k} - k = b$ .*

*Proof.* Using (4), we see that  $S^i(k/N) \equiv (k - ip)/N \pmod{1}$ , which is equivalent to  $NS^i(k/N) \equiv k - ip \pmod{N}$ . We know that the numbers  $p$  and  $N$  are co-prime, thus the mapping  $f_k : \{0, \dots, N - 1\} \rightarrow \{0, \dots, N - 1\}$  given by the congruence  $f_k(i) \equiv k - ip \pmod{N}$  is a bijection. As well,  $f_{\bar{k}}(i) - f_k(i) \equiv \bar{k} - k \pmod{N}$ .

Denote  $m = \min\{p, q\}$  and  $b = \bar{k} - k$ . Consider the following cases:

( $b < 0$ ) We shall see that  $w^{(k)}$  is lexicographically larger than  $w^{(\bar{k})}$ , i.e. if  $i \in \mathbb{N}$  is the first position such that  $w_i^{(k)} \neq w_i^{(\bar{k})}$ , then  $w_i^{(k)} = 1$  and  $w_i^{(\bar{k})} = 0$ . Directly from the definition of amicability, if  $w^{(k)} \propto w^{(\bar{k})}$  and  $w^{(k)} \neq w^{(\bar{k})}$ , then  $w^{(k)}$  is lexicographically smaller than  $w^{(\bar{k})}$ . These two facts make a contradiction.

( $b \in \{0, \dots, m\}$ ) Let  $I_a \subset \{0, \dots, N - 1\}$  be a set of indices  $i$  such that  $w_i^{(k)} = a$  and  $w_i^{(\bar{k})} \neq a$ , for both  $a = 0, 1$ . To show that  $w^{(k)}$  is  $b$ -amicable to  $w^{(\bar{k})}$ , we need to show that  $i \in I_0$  implies  $i + 1 \in I_1$  and  $\#I_0 = \#I_1 = b$ . The fact that  $|w^{(k)}|_0 = |w^{(\bar{k})}|_0$  follows to  $\#I_0 = \#I_1$ .

Let  $i$  be an index such that  $f_k(i) \in [p - b, p)$ , thus  $w_i^{(k)} = 0$ . Then  $f_{\bar{k}}(i) \in [p, p + b)$ , thus  $w_i^{(\bar{k})} = 1$ . This means  $i \in I_0$ . For these  $i$ , we have  $f_k(i + 1) \in [N - b, N)$  and  $f_{\bar{k}}(i + 1) \in [0, b)$ , which means  $i \in I_1$ . There are exactly  $b$  such indices  $i$ .

It remains to show that we covered the whole set  $I_0$ . Suppose  $f_k(i) < p - b$ , then  $f_{\bar{k}}(i) < p$  and  $w_i^{(\bar{k})} = 0$ , which means  $i \notin I_0$ . Suppose  $f_k(i) \geq p$ , then  $w_i^{(k)} = 1$ , which means  $i \notin I_0$ .

( $b \in \{m+1, \dots, N-m-1\}$ ) Let  $i$  be such index that  $f_k(i) = p-1$ .

Suppose  $p \leq q$ . Then  $f_k(i+1) = N-1$ ,  $f_{\bar{k}}(i) = b+p-1$  and  $f_{\bar{k}}(i+1) = b-1$ , which means that  $w_i^{(k)}w_{i+1}^{(k)} = 01$  and  $w_i^{(\bar{k})}w_{i+1}^{(\bar{k})} = 11$ .

Suppose  $p > q$ . Then  $f_k(i+1) = N-1$ ,  $f_{\bar{k}}(i) = b-q-1$  and  $f_{\bar{k}}(i+1) = b-1$ , which means that  $w_i^{(k)}w_{i+1}^{(k)} = 01$  and  $w_i^{(\bar{k})}w_{i+1}^{(\bar{k})} = 00$ .

Both these are in contradiction with  $w^{(k)} \propto w^{(\bar{k})}$ .

( $b \in \{N-m, \dots, N-1\}$ ) We can see that  $N-m = p+q - \min\{p, q\} = \max\{p, q\}$ .

Suppose  $p < q$ . We will show that  $j = 2p$  solves the inequalities

$$\begin{aligned} p \leq j < N, & & p \leq j+b-N < N, \\ p \leq j-p < N, & & 0 \leq j+b-p-N < p. \end{aligned}$$

We have  $2p > p$ ;  $2p < p+q = N$ ;  $2p+b-N \geq 2p+q-N = p$ ;  $2p+b-N < 2p < N$ ;  $2p-p = p \geq p$ ;  $2p-p = p < N$ ;  $2p+b-p-N = p-(N-b) < p$ ;  $2p+b-p-N = b-(N-p) = b-q \geq 0$ .

Let  $i$  be index such that  $f_k(i) = j$ . Then the previous inequalities give  $w_i^{(k)}w_{i+1}^{(k)} = 11$  and  $w_i^{(\bar{k})}w_{i+1}^{(\bar{k})} = 10$ , which is in contradiction with  $w^{(k)} \propto w^{(\bar{k})}$ .

Suppose  $p > q$ . We will show that  $j = \max\{2p-N, N-b\}$  solves the inequalities

$$\begin{aligned} 0 \leq j < p, & & 0 \leq j+b-N < p, \\ p \leq j-p+N < N, & & 0 \leq j+b-p < p. \end{aligned}$$

We have  $j \geq N-b > 0$ , thus  $j > 0$ ;  $2p-N = p+(p-N) < p$  and  $N-b < p$ , thus  $j < p$ ;  $j \geq N-b$ , thus  $j+b-N \geq 0$ ;  $(2p-N)+b-N = p-q-(N-b) < p$  and  $(N-b)+b-N = 0 < p$ , thus  $j+b-N < p$ ;  $j \geq 2p-N$ , thus  $j-p+N \geq p$ ;  $j-p < 0$ , thus  $j-p+N < N$ ;  $j \geq N-b$ , thus  $j+b-p \geq N-b+b-p = q > 0$ ;  $(2p-N)+b-p < 2p-b+b-p = p$  and  $(N-b)+b-p = q < p$ , thus  $j+b-p < p$ .

Let  $i$  be index such that  $f_k(i) = j$ . Then the previous inequalities give  $w_i^{(k)}w_{i+1}^{(k)} = 01$  and  $w_i^{(\bar{k})}w_{i+1}^{(\bar{k})} = 00$ , which is contradiction with  $w^{(k)} \propto w^{(\bar{k})}$ .  $\square$

*Proof of Proposition 11.* Let  $S$  be a 2-interval exchange transformation with the slope  $\varepsilon = p/N$ . Let  $k \in \mathbb{Z}$  and denote  $w^{(k)}$  the word of the length  $N = \|\mathbf{A}\|$  that codes the orbit of the point  $\{k/N\}$  with respect to  $S$ . We know that for every Sturmian morphism  $\varphi$  with  $\mathbf{M}_\varphi = \mathbf{A}$ , there exists

$k \in \{0, \dots, N-1\}$  such that  $\varphi(01) = w^{(k)}$ , we will denote this morphism  $\varphi^{(k)}$ .

Let  $\varphi_{\text{std}}$  be a standard morphism with  $\mathbf{M}_{\varphi_{\text{std}}} = \mathbf{A}$ . Every Sturmian morphism  $\varphi^{(k)}$  is a right conjugate to  $\varphi_{\text{std}}$ , which means that there exist words  $v, v' \in \{0, 1\}^*$  such that

$$\varphi(aa') = v01v' \quad \text{and} \quad \varphi(a'a) = v10v',$$

where letters  $a, a'$  satisfy  $aa' = 01$  for  $\det \mathbf{A} = +1$  and  $aa' = 10$  for  $\det \mathbf{A} = -1$ . This gives that  $\varphi(aa')$  is 1-amicable to  $\varphi(a'a)$ .

Morphism  $\varphi^{(k)}$  is  $b$ -amicable to  $\varphi^{(\bar{k})}$  if and only if the following conditions are satisfied:

1.  $\varphi^{(k)}(01)$  is  $b$ -amicable to  $\varphi^{(\bar{k})}(10)$ ;
2.  $\varphi^{(k)}(01)$  is amicable to  $\varphi^{(\bar{k})}(01)$ ;
3. Parikh vectors satisfy  $\mathbf{M}_{\varphi^{(k)}(0)} = \mathbf{M}_{\varphi^{(\bar{k})}(0)}$ .

The 2nd and 3rd conditions assures that  $\varphi^{(k)}(0) \propto \varphi^{(\bar{k})}(0)$  and  $\varphi^{(k)}(1) \propto \varphi^{(\bar{k})}(1)$ .

Let us discuss the cases  $\det \mathbf{A} = +1$  and  $\det \mathbf{A} = -1$ .

( $\det \mathbf{A} = +1$ ) We know that  $\varphi^{(k)}(01)$  is 1-amicable to  $\varphi^{(k)}(10)$ , which implies  $\varphi^{(k)}(10) = w^{(k+1)}$ . This excludes  $k = N-1$ .

The 3rd condition is immediately satisfied by  $\mathbf{M}_{\varphi^{(k)}} = \mathbf{M}_{\varphi^{(\bar{k})}}$ . To satisfy the 1st condition, we need  $(\bar{k} + 1) - k = b$ . To satisfy the 2nd condition, we need  $0 \leq \bar{k} - k \leq \min\{p, q\}$ . These facts gives  $0 \leq k \leq \bar{k} \leq N-2$  and  $1 \leq b \leq \min\{p, q\}$ . For each such  $b$ , we have exactly  $N - b$  pairs of such indices  $(k, \bar{k})$ .

( $\det \mathbf{A} = -1$ ) We know that  $\varphi^{(k)}(10)$  is 1-amicable to  $\varphi^{(k)}(01)$ , which implies  $\varphi^{(k)}(10) = w^{(k-1)}$ . This excludes  $k = 0$ .

The 3rd condition is immediately satisfied by  $\mathbf{M}_{\varphi^{(k)}} = \mathbf{M}_{\varphi^{(\bar{k})}}$ . To satisfy the 1st condition, we need  $(\bar{k} - 1) - k = b$ . To satisfy the 2nd condition, we need  $0 \leq \bar{k} - k \leq \min\{p, q\}$ . These facts gives  $1 \leq k \leq \bar{k} \leq N-1$  and  $0 \leq b \leq \min\{p, q\} - 1$ . For each such  $b$ , we have exactly  $N - b - 2$  pairs of such indices  $(k, \bar{k})$ .  $\square$

*Proof of Theorem 1.* The formula (2) can be obtained by summation of numbers  $c_{\mathbf{A}}(b)$  from the previous proposition.  $\square$

### 3.3 Matrices of ternarizations

To each pair of amicable Sturmian morphisms, an incidence matrix of its ternarization is assigned. We now fully describe which matrices from  $\mathbb{N}^{3 \times 3}$  are matrices of ternarizations.

**Theorem 12.** *A matrix  $\mathbf{B} \in \mathbb{N}^{3 \times 3}$  is the incidence matrix of the ternarization of a pair of amicable Sturmian morphisms if and only if there exist matrix  $\mathbf{A} = \begin{pmatrix} p_0 & q_0 \\ p_1 & q_1 \end{pmatrix} \in \mathbb{N}^{2 \times 2}$  with  $\det \mathbf{A} = \Delta = \pm 1$  and numbers  $b_0, b_1 \in \mathbb{N}$  such that*

- (a)  $\left| \frac{b_0(p_1+q_1)-b_1(p_0+q_0)}{p_0+q_0+p_1+q_1} \right| < 1,$
- (b)  $\frac{1-\Delta}{2} \leq b_0 + b_1 \leq \min\{p_0 + p_1, q_0 + q_1\} - \frac{\Delta+1}{2},$
- (c)  $\mathbf{B} = \mathbf{P} \begin{pmatrix} \mathbf{A} & b_0 \\ & b_1 \\ 0 & 0 & \Delta \end{pmatrix} \mathbf{P}^{-1},$  where  $\mathbf{P} = \begin{pmatrix} 1 & 0 & 0 \\ 1 & 1 & 1 \\ 0 & 1 & 0 \end{pmatrix}.$

*Proof of the implication ( $\Rightarrow$ ).* Let us denote  $p = p_0 + p_1$ ,  $q = q_0 + q_1$ ,  $N = p + q$  and  $b = b_0 + b_1 + \Delta$ . Then we can see that condition (c) gives

$$\mathbf{B} = \begin{pmatrix} p_0 - b_0 & b_0 & q_0 - b_0 \\ p - b & b & q - b \\ p_1 - b_1 & b_1 & q_1 - b_1 \end{pmatrix}. \quad (7)$$

The fact that (c) is necessary for  $\mathbf{B}$  to be an incidence matrix of a ternarization is shown in [3]. Condition (b) is necessary according to Proposition 11, so we only need to show that (a) is satisfied for the matrix of the ternarization  $\eta = \text{ter}(\varphi, \psi)$  of a pair of amicable Sturmian morphisms  $\varphi \propto \psi$ .

We can see that  $\mathbf{A} = \begin{pmatrix} p_0 & q_0 \\ p_1 & q_1 \end{pmatrix}$  is necessarily an incidence matrix of both  $\varphi$  and  $\psi$ . Let  $S$  be a 2-interval exchange transformation with rational slope  $\varepsilon = p/N$ . Then there exist  $k, \bar{k} \in \{0, \dots, N-2\}$  such that  $\varphi(01)$ ,  $\psi(01)$  code transformation  $S$  with start points  $x_0 = k/N$ ,  $\bar{x}_0 = \bar{k}/N$ ; moreover,  $\bar{k} - k = b - \Delta$ . We need to determine the value of  $b_0 = |\text{ter}(\varphi(0), \psi(0))|_B$ . The number  $b_0$  is equal to the number of indices  $i \in \{0, 1, \dots, p_0 + q_0 - 1\}$  such that  $S^i x_0 \in [(p-b+\Delta)/N, p/N)$ , because for exactly these  $i$ , we have  $S^i x_0 < p/N \leq S^{i+1} x_0$ .

Let  $X = \{x_0 - ip/N \mid i \in \mathbb{N}, 0 \leq i < p_0 + q_0\}$ . Put  $p' = p + \Delta/k$ , and let  $Y = \{x_0 - ip'/N \mid i \in \mathbb{N}, 0 \leq i < p_0 + q_0\}$ . We can see that  $0 \leq \Delta((x_0 - ip/N) - (x_0 - ip'/N)) = i/kN < 1/N$ . Thus  $x_0 - ip/N \in [\frac{p-b+\Delta}{N}, \frac{p}{N})$  if and only if

$$x_0 - ip'/N \in \begin{cases} (\frac{p-b}{N}, \frac{p-1}{N}] & \text{in the case } \Delta = +1, \\ [\frac{p-b-1}{N}, \frac{p}{N}) & \text{in the case } \Delta = -1. \end{cases} \quad (8)$$

In both cases, the length of the interval is  $\frac{b-\Delta}{N}$ . From  $\det \mathbf{A} = \Delta$ , it is easy to see that  $p'/N = p_0/(p_0 + q_0)$ . Because  $p_0$  is co-prime to  $p_0 + q_0$ , we get

$\{i p_0 / (p_0 + q_0) \mid i \in \mathbb{N}, 0 \leq i < p_0 + q_0\} = \{i / (p_0 + q_0) \mid i \in \mathbb{N}, 0 \leq i < p_0 + q_0\}$ . But this means that the set  $Y$  is uniformly distributed on the interval  $[0, 1)$ , which gives

$$b_0 = \# \left( X \cap \left[ \frac{p-b+\Delta}{N}, \frac{p}{N} \right) \right) \in \{\lfloor \beta \rfloor, \lceil \beta \rceil\},$$

where  $\beta = (p_0 + q_0) \frac{b-\Delta}{N}$  is number of elements of  $Y$  multiplied by the length of the interval in (8). Together we get

$$|\beta - b_0| < 1, \quad (9)$$

which is equivalent to condition (a).  $\square$

The proof of the other implication is divided into several lemmas.

**Lemma 2.** Let  $\mathbf{A} = \begin{pmatrix} p_0 & q_0 \\ p_1 & q_1 \end{pmatrix} \in \mathbb{N}^{2 \times 2}$  with  $\det \mathbf{A} = \Delta = \pm 1$ , let  $b \in \mathbb{N}$  with  $\frac{1+\Delta}{2} \leq b \leq \min\{p_0 + p_1, q_0 + q_1\} - \frac{1-\Delta}{2}$ .

Denote  $N = \|\mathbf{A}\|$  and  $p = p_0 + p_1$  integers,  $I = \left[ \frac{p-b+\Delta}{N}, \frac{p}{N} \right)$  an interval,  $X_k = \{k/N, S\{k/N\}, S^2\{k/N\}, \dots, S^{p_0+q_0-1}\{k/N\}\}$  a set of numbers for any  $k \in \mathbb{Z}$ , where  $S$  is the 2-interval exchange with the slope  $\varepsilon = p/N$ , and denote  $\beta = \frac{p_0+q_0}{N}(b-\Delta)$ .

Then for all  $b_0 \in \{\lfloor \beta \rfloor, \lceil \beta \rceil\}$  such that

$$b_0 \leq \min\{p_0, q_0\} \quad \text{and} \quad b - \Delta - b_0 \leq \min\{p_1, q_1\}, \quad (10)$$

there exist  $k', k'' \in \{0, \dots, N-1\}$ ,  $k' \neq k''$  such that

$$\#(X_{k'} \cap I) = \#(X_{k''} \cap I) = b_0. \quad (11)$$

*Proof.* Denote  $r(k) = \#(X_k \cap I)$  for  $k \in \mathbb{Z}$ . We can see that  $\sum_{k=0}^{N-1} r(k) = (b-\Delta)(p_0+q_0)$ . According to (9), we know that  $r(k) \in \{\lfloor \beta \rfloor, \lceil \beta \rceil\}$  for all  $k \in \mathbb{Z}$ . Let

$$\begin{aligned} C_L &= \#\{k \in \{0, \dots, N-1\} \mid r(k) = \lfloor \beta \rfloor\}, \\ C_U &= \#\{k \in \{0, \dots, N-1\} \mid r(k) = \lceil \beta \rceil\}. \end{aligned}$$

We will proof the lemma by contradiction. Suppose  $C_L$  or  $C_U \in \{0, 1\}$ . The numbers satisfy equations

$$C_L \lfloor \beta \rfloor + C_U \lceil \beta \rceil = N\beta \quad \text{and} \quad C_L + C_U = N.$$

If  $C_L = 0$  or  $C_U = 0$ , necessarily  $\beta \in \mathbb{N}$  and (11) is satisfied for all  $k \in \mathbb{Z}$ . Otherwise, there is a unique solution

$$C_L = N\{-\beta\} \quad \text{and} \quad C_U = N\{\beta\}.$$



Using relation  $p_0N - (p_0 + p_1)(p_0 + q_0) = \Delta$ , we get

$$\begin{aligned} C_U &\equiv (p_0 + q_0)(b - \Delta) \pmod{N} \\ b - \Delta &\equiv -\Delta(p_0 + p_1)C_U \pmod{N}. \end{aligned}$$

Suppose  $C_U = 1$  or  $C_L = 1$ , i.e.  $C_U \equiv \pm 1$ . Then  $b = (p_0 + p_1) + \Delta$  or  $b = (q_0 + q_1) + \Delta$ . For  $\Delta = +1$ , this is in contradiction with the conditions. For  $\Delta = -1$ , discuss the following two cases.

1. Suppose  $b = (p_0 + p_1) + \Delta$ . This means  $C_U = 1$ . But then  $b_0 = \lceil \beta \rceil = \lceil \frac{p_0N - \Delta}{N} \rceil = p_0 + 1$  does not satisfy condition (10) of the lemma.
2. Suppose  $b = (q_0 + q_1) + \Delta$ . That means  $C_L = 1$ . But then  $b_0 = \lfloor \beta \rfloor = q_0 - 1$  and  $b - \Delta - b_0 = q_1 + 1$  again does not satisfy (10).  $\square$

**Lemma 3.** Let  $\mathbf{A} = \begin{pmatrix} p_0 & q_0 \\ p_1 & q_1 \end{pmatrix} \in \mathbb{N}^{2 \times 2}$  with  $\det \mathbf{A} = \Delta = \pm 1$ , let  $b \in \mathbb{N}$  with  $\frac{1+\Delta}{2} \leq b \leq \min\{p_0 + p_1, q_0 + q_1\} - \frac{1-\Delta}{2}$ .

Denote  $N = \|\mathbf{A}\|$  and  $p = p_0 + p_1$  integers,  $I = [\frac{p-b+\Delta}{N}, \frac{p}{N})$  an interval,  $X_k = \{\{k/N\}, S\{k/N\}, S^2\{k/N\}, \dots, S^{p_0+q_0-1}\{k/N\}\}$  a set of numbers for any  $k \in \mathbb{Z}$ , where  $S$  is the 2-interval exchange with slope  $\varepsilon = p/N$ , and denote  $\beta = \frac{p_0+q_0}{N}(b - \Delta)$ . Define morphisms  $\varphi_k$  for  $k \in \mathbb{Z}$  in the following way:

- the word  $\varphi_k(0)$  codes  $\{k/N\}, S\{k/N\}, \dots, S^{p_0+q_0-1}\{k/N\}$ ;
- the word  $\varphi_k(1)$  codes  $S^{p_0+q_0}\{k/N\}, \dots, S^{N-1}\{k/N\}$ .

Let  $k_0 \in \mathbb{Z}$  such that  $\#(X_{k_0} \cap I) = \#(X_{k_0-p} \cap I)$ . Then

$$\varphi_{k_0} \propto \varphi_{k_0+b-\Delta} \quad \text{or} \quad \varphi_{k_0-p} \propto \varphi_{k_0-p+b-\Delta},$$

and the number of  $B$ 's in the ternarization of the images of the letter  $0$  is  $\#(X_{k_0} \cap I)$ .

*Proof.* Let us take the orbit

$$\{k_0/N\}, S\{k_0/N\}, \dots, S^{p_0+q_0}\{k_0/N\}. \quad (12)$$

Let  $t^{(k)}$  be a word of the length  $p_0 + q_0$  that codes orbit of transformation  $S$  to the alphabet  $\{0, 0', 1, 1'\}$  with a different code than (3):

$$t_i^{(k)} = \begin{cases} 0 & \text{if } S^i\{k/N\} \in [0, \frac{p-b+\Delta}{N}), \\ 0' & \text{if } S^i\{k/N\} \in [\frac{p-b+\Delta}{N}, \frac{p}{N}) = I, \\ 1 & \text{if } S^i\{k/N\} \in [\frac{p}{N}, \frac{N-b+\Delta}{N}), \\ 1' & \text{if } S^i\{k/N\} \in [\frac{p-b+\Delta}{N}, 1). \end{cases} \quad (13)$$

From definition of  $S$ , we see that  $t_i^{(k)} = 0' \Leftrightarrow t_{i+1}^{(k)} = 1'$ . Define two morphisms  $\tau, \tau' : \{0, 0', 1, 1'\}^* \rightarrow \{0, 1\}^*$  as

$$\begin{aligned} \tau(0) &= 0, & \tau(0') &= 0, & \tau(1) &= 1, & \tau(1') &= 1, \\ \tau'(0) &= 0, & \tau'(0') &= 1, & \tau'(1) &= 1, & \tau'(1') &= 0. \end{aligned}$$

If  $t^{(k)}$  does not start with  $1'$  and does not end with  $0'$ , then the word  $\varphi_k(0) = \tau(t^{(k)})$  is  $|t^{(k)}|_{0'}$ -amicable to  $\tau'(t^{(k)}) = \varphi_{k+b-\Delta}(0)$ . Moreover,  $|t^{(k)}|_{0'} = \#(X_k \cap I)$ .

We know that  $S\{k_0/N\} = \{(k_0 - p)/N\}$ , which means that there exist letters  $a, b \in \{0, 0', 1, 1'\}$  such that  $t^{(k_0)}a = bt^{(k_0-p)}$  and  $a = 0' \Leftrightarrow b = 0'$ , because  $|t^{(k_0)}|_{0'} = |t^{(k_0-p)}|_{0'}$ . If  $a = 0'$  then  $t^{(k_0)}$  does not end with  $0'$ , because in that case  $b = 1'$ . If  $a \neq 0'$  then  $t^{(k_0-p)}$  does not start with  $1'$  and does not end with  $0'$ . Putting these facts together with facts from the proof of Proposition 11 we get the statement.  $\square$

**Lemma 4.** Let  $\mathbf{A} = \begin{pmatrix} p_0 & q_0 \\ p_1 & q_1 \end{pmatrix} \in \mathbb{N}^{2 \times 2}$  with  $\det \mathbf{A} = \Delta = \pm 1$ , let  $b \in \mathbb{N}$  with  $\frac{1+\Delta}{2} \leq b \leq \min\{p_0 + p_1, q_0 + q_1\} - \frac{1-\Delta}{2}$ .

Denote  $N = \|\mathbf{A}\|$ ,  $q = q_0 + q_1$  and  $p = p_0 + p_1$  integers,  $I = [\frac{p-b+\Delta}{N}, \frac{p}{N})$  an interval,  $X_k = \{\{k/N\}, S\{k/N\}, S^2\{k/N\}, \dots, S^{p_0+q_0-1}\{k/N\}\}$  a set of numbers for any  $k \in \mathbb{Z}$ , where  $S$  is the 2-interval exchange with the slope  $\varepsilon = p/N$ , and denote  $\beta = \frac{p_0+q_0}{N}(b - \Delta)$ .

Let  $k_0 \in \mathbb{Z}$  be a number such that if  $\Delta = -1$  and  $b = \min\{p, q\} - 1$  then

$$k_0 \not\equiv \begin{cases} -1 \pmod{N} & \text{in the case } p > q, \\ p - b - 1 \pmod{N} & \text{in the case } p < q. \end{cases} \quad (14)$$

Then

$$\#(X_{k_0} \cap I) = \#(X_{k_0+p} \cap I) \quad \text{or} \quad \#(X_{k_0} \cap I) = \#(X_{k_0-p} \cap I).$$

*Proof.* Define the words  $t^{(k)}$  by (13) in the same way as in the previous proof. Denote  $\ell = p_0 + q_0$ . Then we know that there exist letters  $a_0, \dots, a_{\ell+1} \in \{0, 0', 1, 1'\}$  such that

$$\begin{aligned} t^{(k_0+p)} &= a_0 a_1 a_2 \cdots a_{\ell-1}, \\ t^{(k_0)} &= a_1 a_2 \cdots a_{\ell-1} a_{\ell}, \\ t^{(k_0-p)} &= a_2 \cdots a_{\ell-1} a_{\ell} a_{\ell+1}. \end{aligned}$$

Remind that  $\#(X_{k_0+p} \cap I) = |t^{(k)}|_{0'}$ . The proof will be done by contradiction. Suppose that  $|t^{(k_0+p)}|_{0'} \neq |t^{(k_0)}|_{0'} \neq |t^{(k_0-p)}|_{0'}$ . There are only two possible values of these numbers, thus  $|t^{(k_0+p)}|_{0'} = |t^{(k_0-p)}|_{0'}$ . This together gives either  $a_0 = a_{\ell+1} = 0'$  or  $a_1 = a_{\ell} = 0'$ . It means that there exist

$\xi \in I = \left[\frac{p-b+\Delta}{N}, \frac{p}{N}\right)$  and  $\omega = \pm 1$  such that  $S^{\ell+\omega}\xi \in I$ . We can take  $\xi \in \frac{1}{N}\mathbb{Z}$ . Since  $\ell p = p_0 N - \Delta$ , we have

$$S^{\ell+\omega}\xi \equiv \xi - \frac{(\ell + \omega)p}{N} \equiv \frac{\omega p - \Delta}{N} \pmod{1}.$$

This gives

$$\begin{aligned} S^{\ell+\omega}\xi - \xi &= \frac{p - \omega\Delta}{N} \\ \text{or } S^{\ell+\omega}\xi - \xi &= \frac{p - \omega\Delta}{N} - 1 = -\frac{q + \omega\Delta}{N}. \end{aligned}$$

This enforces  $b - 1 - \Delta \geq \min\{p, q\} - 1$  for the interval  $I$  to be large enough to contain both  $\xi$  and  $S^{\ell+\omega}\xi$ .

For  $\Delta = +1$ , this is in contradiction with  $b \leq \min\{p, q\}$ .

For  $\Delta = -1$  we get only one admissible  $b = \min\{p, q\} - 1$ . If  $p = \min\{p, q\}$ , it gives  $\omega = -1$  and  $\xi = \frac{p-b-1}{N}$ , which implies  $k_0 \equiv p - b - 1 \pmod{N}$ . If  $q = \min\{p, q\}$ , it gives  $\omega = +1$  and  $\xi = \frac{p-1}{N}$ , which implies  $k_0 \equiv -1 \pmod{N}$ . Both these cases are in contradiction with (14).  $\square$

*Proof of the implication ( $\Leftarrow$ ).* From [3], the incidence matrix of the ternarization  $\text{ter}(\varphi, \psi)$  is fully described by the matrix  $\mathbf{A}$  and numbers  $b_0$  and  $b = b_0 + b_1 + \Delta$ . The condition (a) is equivalent to (9) and it gives at most two values of  $b_0$ . If  $\beta \in \mathbb{N}$ , there is nothing to do as we have at least one pair of  $b$ -amicable morphisms  $\varphi \propto \psi$  for  $\mathbf{A}$ , and its incidence matrix satisfies all three conditions.

For  $\beta \notin \mathbb{N}$ , we want to show that for both  $b_0 \in \{[\beta], \lceil\beta\rceil\}$  there exist  $\varphi \propto \psi$  with  $|\text{ter}(\varphi(0), \psi(0))|_B = b_0$ . Because the elements of the matrix  $\mathbf{B}$  are non-negative, the condition (10) of Lemma 2 is satisfied and we have two different  $k', k''$ . At least one of them satisfies (14). Lemma 4 then provides  $k_0$  satisfying the conditions of Lemma 3 that gives a pair of amicable Sturmian morphisms, ternarization of which has the incidence matrix  $\mathbf{B}$ .  $\square$

## 4 Conclusions and open problems

Matrices of 3iet-preserving morphisms were studied in [3]. The authors give a necessary condition on  $\mathbf{B} \in \mathbb{N}^{3 \times 3}$  to be an incidence matrix of a 3iet-preserving morphism:

$$\mathbf{BEB}^\top = \pm \mathbf{E}, \quad \text{where } \mathbf{E} = \begin{pmatrix} 0 & 1 & 1 \\ -1 & 0 & 1 \\ -1 & -1 & 0 \end{pmatrix}.$$

However, this condition is not sufficient. In our contribution, we study 3iet-preserving morphisms  $\eta = \text{ter}(\varphi, \psi)$  arising from pairs of amicable Sturmian

morphisms  $\varphi \propto \psi$ . Our Theorem 12 gives sufficient and necessary condition for any matrix  $\mathbf{B} \in \mathbb{N}^{3 \times 3}$  to satisfy  $\mathbf{B} = \mathbf{M}_\eta$  for some ternarization  $\eta = \text{ter}(\varphi, \psi)$ .

It remains to answer the question about the role of the monoid

$$\mathcal{M}_{\text{ter}} = \{ \text{ter}(\varphi, \psi) \mid \varphi, \psi \text{ amicable morphisms} \}$$

in the whole monoid  $\mathcal{M}_{3\text{iet}}$  of all 3iet-preserving morphisms.

It seems that using similar proof as [2] for Theorem 5 we can proof the following statement.

**Conjecture.** *Let  $\eta \in \mathcal{M}_{3\text{iet}}$ . Then there exists  $i \in \{0, 1, 2, 3\}$  such that  $\eta \circ \xi_i \in \mathcal{M}_{\text{ter}}$ , where  $\xi_0, \dots, \xi_3$  are 3iet-preserving morphisms,*

$$\begin{array}{llll} \xi_0(A) = A, & \xi_1(A) = C, & \xi_2(A) = B, & \xi_3(A) = B, \\ \xi_0(B) = B, & \xi_1(B) = B, & \xi_2(B) = ACA, & \xi_3(B) = CAC, \\ \xi_0(C) = C, & \xi_1(C) = A, & \xi_2(C) = A, & \xi_3(C) = C. \end{array}$$

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